

7 Binary relations

Definition 7.1. A *binary relation* on a set A is a subset $R \subseteq A \times A$. We usually write aRb instead of writing $\langle a, b \rangle \in R$.

Example 7.2. 1. The order relation on \mathbb{N} is given by

$$\{\langle n, m \rangle \mid n, m \in \mathbb{N}, n < m\}.$$

2. The division relation D on $\mathbb{N} \setminus \{0\}$ is given by

$$D = \{\langle n, m \rangle \mid n, m \in \mathbb{N}, n \text{ divides } m\}.$$

Observation Thus $\mathcal{P}(\mathbb{N} \times \mathbb{N})$ is the collection of all binary relations on \mathbb{N} . Clearly $\mathcal{P}(\mathbb{N} \times \mathbb{N}) \sim \mathcal{P}(\mathbb{N})$ and so $\mathcal{P}(\mathbb{N} \times \mathbb{N})$ is uncountable.

Definition 7.3. Let R be a binary relation on A .

1. R is *reflexive* iff xRx for all $x \in A$.
2. R is *symmetric* iff xRy implies yRx for all $x, y \in A$.
3. R is *transitive* iff xRy and yRz implies xRz for all $x, y, z \in A$.

R is an *equivalence relation* iff R is reflexive, symmetric, and transitive.

Example 7.4. Define the relation R on \mathbb{Z} by

$$aRb \text{ iff } 3 \mid a - b.$$

Proposition 7.5. R is an equivalence relation.

Exercise 7.6. Let $A = \{\langle a, b \rangle \mid a, b \in \mathbb{Z}, b \neq 0\}$. Define the relation S on A by

$$\langle a, b \rangle S \langle c, d \rangle \text{ iff } ad - bc = 0.$$

Prove that S is an equivalence relation.

Definition 7.7. Let R be an equivalence relation on A . For each $x \in A$, the *equivalence class* of x is

$$[x] = \{y \in A \mid xRy\}.$$

Example 7.4 Cont. The distinct equivalence classes are

$$[0] = \{\dots, -6, -3, 0, 3, 6, \dots\}$$

$$[1] = \{\dots, -5, -2, 1, 4, 7, \dots\}$$

$$[2] = \{\dots, -4, -1, 2, 5, 8, \dots\}$$

Definition 7.8. Let A be a nonempty set. Then $\{B_i \mid i \in I\}$ is a *partition* of A iff the following conditions hold:

1. $\emptyset \neq B_i$ for all $i \in I$.
2. If $i \neq j \in I$, then $B_i \cap B_j = \emptyset$.
3. $A = \bigcup_{i \in I} B_i$.

Theorem 7.9. *Let R be an equivalence relation on A .*

1. *If $a \in A$ then $a \in [a]$.*
2. *If $a, b \in A$ and $[a] \cap [b] \neq \emptyset$, then $[a] = [b]$.*

Hence the set of distinct equivalence classes forms a partition of A .

Proof. 1. Let $a \in A$. Since R is reflexive, aRa and so $a \in [a]$.

2. Suppose that $c \in [a] \cap [b]$. Then aRc and bRc . Since R is symmetric, cRb . Since R is transitive, aRb . We claim that $[b] \subseteq [a]$. To see this, suppose that $d \in [b]$. Then bRd . Since aRb and bRd , it follows that aRd . Thus $d \in [a]$. Similarly, $[a] \subseteq [b]$ and so $[a] = [b]$. □

Theorem 7.10. *Let $\{B_i \mid i \in I\}$ be a partition of A . Define a binary relation R on A by*

$$aRb \text{ iff there exists } i \in I \text{ such that } a, b \in B_i.$$

Then R is an equivalence relation whose equivalence classes are precisely $\{B_i \mid i \in I\}$. □

Example 7.11. How many equivalence relations can be defined on $A = \{1, 2, 3\}$?

Sol'n This is the same as asking how many partitions of A exist.

$$\begin{aligned} & \{\{1, 2, 3\}\}, \\ & \{\{1, 2\}, \{3\}\}, \{\{1, 3\}, \{2\}\}, \{\{2, 3\}, \{1\}\}, \\ & \{\{1\}, \{2\}, \{3\}\} \end{aligned}$$

Hence there are 5 equivalence relations on $\{1, 2, 3\}$.

Exercise 7.12. How many equivalence relations can be defined on $A = \{1, 2, 3, 4\}$?

Challenge Let $\text{EQ}(\mathbb{N})$ be the collection of equivalence relations on \mathbb{N} . Prove that $\text{EQ}(\mathbb{N}) \sim \mathcal{P}(\mathbb{N})$.

8 Linear orders

Definition 8.1. Let R be a binary relation on A .

1. R is *irreflexive* iff $\langle a, a \rangle \notin R$ for all $a \in A$.
2. R satisfies the *trichotomy property* iff for all $a, b \in A$, exactly one of the following holds:

$$aRb, a = b, bRa.$$

$\langle A, R \rangle$ is a *linear order* iff R is irreflexive, transitive, and satisfies the trichotomy property.

Example 8.2. Each of the following are linear orders.

1. $\langle \mathbb{N}, < \rangle$
2. $\langle \mathbb{N}, > \rangle$
3. $\langle \mathbb{Z}, < \rangle$
4. $\langle \mathbb{Q}, < \rangle$
5. $\langle \mathbb{R}, < \rangle$

Definition 8.3. Let R be a binary relation on A . Then $\langle A, R \rangle$ is a *partial order* iff R is irreflexive and transitive.

Example 8.4. Each of the follow are partial orders, but *not* linear orders.

1. Let A be any nonempty set containing at least two elements. Then $\langle \mathcal{P}(A), \subset \rangle$ is a partial order.
2. Let D be the divisibility relation on $\mathbb{N}^+ = \mathbb{N} \setminus \{0\}$. Then $\langle \mathbb{N}^+, D \rangle$ is a partial order.

Definition 8.5. Let $\langle A, < \rangle$ and $\langle B, < \rangle$ be partial orders. A map $f: A \rightarrow B$ is an *isomorphism* iff the following conditions are satisfied.

1. f is a bijection
2. For all $x, y \in A$, $x < y$ iff $f(x) < f(y)$.

In this case, we say that $\langle A, < \rangle$ and $\langle B, < \rangle$ are isomorphic and write $\langle A, < \rangle \cong \langle B, < \rangle$.

Example 8.6. $\langle \mathbb{Z}, < \rangle \cong \langle \mathbb{Z}, > \rangle$

Proof. Let $f: \mathbb{Z} \rightarrow \mathbb{Z}$ be the map defined by $f(x) = -x$. Clearly f is a bijection. Also, if $x, y \in \mathbb{Z}$, then $x < y$
iff $-x > -y$
iff $f(x) > f(y)$.

Thus f is an isomorphism. □

Example 8.7. $\langle \mathbb{N}, < \rangle \not\cong \langle \mathbb{Z}, < \rangle$.

Proof. Suppose that $f: \mathbb{N} \rightarrow \mathbb{Z}$ is an isomorphism. Let $f(0) = z$. Since f is a bijection, there exists $n \in \mathbb{N}$ such that $f(n) = z - 1$. But then $n > 0$ and $f(n) < f(0)$, which is a contradiction. □

Exercise 8.8. Prove that $\langle \mathbb{Z}, < \rangle \not\cong \langle \mathbb{Q}, < \rangle$.

Example 8.9. $\langle \mathbb{Q}, < \rangle \not\cong \langle \mathbb{R}, < \rangle$.

Proof. Since \mathbb{Q} is countable and \mathbb{R} is uncountable, there does not exist a bijection $f: \mathbb{Q} \rightarrow \mathbb{R}$. Hence there does not exist an isomorphism $f: \mathbb{Q} \rightarrow \mathbb{R}$. □

Example 8.10. $\langle \mathbb{R}, < \rangle \not\cong \langle \mathbb{R} \setminus \{0\}, < \rangle$.

Proof. Suppose that $f: \mathbb{R} \setminus \{0\} \rightarrow \mathbb{R}$ is an isomorphism. For each $n \geq 1$, let $r_n = f(1/n)$. Then

$$r_1 > r_2 > \dots > r_n > \dots > f(-1).$$

Let s be the greatest lower bound of $\{r_n \mid n \geq 1\}$. Then there exists $t \in \mathbb{R} \setminus \{0\}$ such that $f(t) = s$. Clearly $t < 0$. Hence $f(t/2) > s$. But then there exists $n \geq 1$ such that $r_n < f(t/2)$. But this means that $t/2 < 1/n$ and $f(t/2) > f(1/n)$, which is a contradiction. □

Question 8.11. Is $\langle \mathbb{Q}, < \rangle \cong \langle \mathbb{Q} \setminus \{0\}, < \rangle$?

Definition 8.12. For each prime p ,
 $\mathbb{Z}[1/p] = \{a/p^n \mid a \in \mathbb{Z}, n \in \mathbb{N}\}$.

Question 8.13. Is $\langle \mathbb{Z}[1/2], < \rangle \cong \langle \mathbb{Z}[1/3], < \rangle$?

Definition 8.14. A linear order $\langle D, < \rangle$ is a *dense linear order without endpoints* or DLO iff the following conditions hold.

1. For all $a, b \in D$, if $a < b$, then there exists $c \in D$ such that $a < c < b$.
2. For all $a \in D$, there exists $b \in D$ such that $a < b$.
3. For all $a \in D$, there exists $b \in D$ such that $b < a$.

Example 8.15. The following are DLOs.

1. $\langle \mathbb{Q}, < \rangle$
2. $\langle \mathbb{R}, < \rangle$
3. $\langle \mathbb{Q} \setminus \{0\}, < \rangle$
4. $\langle \mathbb{R} \setminus \{0\}, < \rangle$

Theorem 8.16. For each prime p , $\langle \mathbb{Z}[1/p], < \rangle$ is a DLO.

Proof. Clearly $\langle \mathbb{Z}[1/p], < \rangle$ linear order without endpoints. Hence it is enough to show that $\mathbb{Z}[1/p]$ is dense. Suppose that $a, b \in \mathbb{Z}[1/p]$. Then there exists $c, d \in \mathbb{Z}$ and $n \in \mathbb{N}$ such that $a = c/p^n$ and $b = d/p^n$. Clearly $a < a + (1/p^n) \leq b$. Consider

$$r = \frac{c}{p^n} + \frac{1}{p^n} = \frac{cp+1}{p^{n+1}} \in \mathbb{Z}[1/p].$$

Then $a < r < b$. □

Theorem 8.17. If $\langle A, < \rangle$ and $\langle B, < \rangle$ are countable DLOs then $\langle A, < \rangle \cong \langle B, < \rangle$.

Corollary 8.18. $\langle \mathbb{Q}, < \rangle \cong \langle \mathbb{Q} \setminus \{0\}, < \rangle$. □

Corollary 8.19. $\langle \mathbb{Z}[1/2], < \rangle \cong \langle \mathbb{Z}[1/3], < \rangle$. □

Corollary 8.20. If p is any prime, then $\langle \mathbb{Z}[1/p], < \rangle \cong \langle \mathbb{Q}, < \rangle$. □

Proof of Theorem 8.17. Let $A = \{a_n \mid n \in \mathbb{N}\}$ and $B = \{b_n \mid n \in \mathbb{N}\}$. First define $A_0 = \{a_0\}$ and $B_0 = \{b_0\}$ and let $f_0: A_0 \rightarrow B_0$ be the map defined by $f_0(a_0) = b_0$.

Now suppose inductively that we have defined a function $f_n: A_n \rightarrow B_n$ such that the following conditions are satisfied.

1. $\{a_0, \dots, a_n\} \subseteq A_n \subseteq A$.
2. $\{b_0, \dots, b_n\} \subseteq B_n \subseteq B$.
3. $f_n: A_n \rightarrow B_n$ is an order preserving bijection.

We now extend f_n to a suitable function f_{n+1} .

Step 1 If $a_{n+1} \in A_n$, then let $A'_n = A_n$, $B'_n = B_n$, and $f'_n = f_n$. Otherwise, suppose for example that

$$c_0 < c_1 < \dots < c_i < a_{n+1} < c_{i+1} < \dots < c_m$$

where $A_n = \{c_0, \dots, c_m\}$. Choose some element $b \in B$ such that $f_n(c_i) < b < f_n(c_{i+1})$ and define

$$\begin{aligned} A'_n &= A_n \cup \{a_{n+1}\} \\ B'_n &= B_n \cup \{b\} \\ f'_n &= f_n \cup \{(a_{n+1}, b)\} \end{aligned}$$

Step 2 If $b_{n+1} \in B'_n$, then let $A_{n+1} = A'_n$, $B_{n+1} = B'_n$, and $f_{n+1} = f'_n$. Otherwise, suppose for example that

$$d_0 < d_1 < \dots < d_j < b_{n+1} < d_{j+1} < \dots < d_t$$

where $B'_n = \{d_0, \dots, d_t\}$. Choose some element $a \in A$ such that $(f'_n)^{-1}(d_j) < a < (f'_n)^{-1}(d_{j+1})$ and define

$$\begin{aligned} A_{n+1} &= A'_n \cup \{a\} \\ B_{n+1} &= B'_n \cup \{b_{n+1}\} \\ f_{n+1} &= f'_n \cup \{\langle a, b_{n+1} \rangle\}. \end{aligned}$$

Finally, let $f = \bigcup_{n \geq 0} f_n$. Then $f: A \rightarrow B$ is an isomorphism. \square

9 Extensions

Definition 9.1. Suppose that R, S are binary relations on A . Then S *extends* R iff $R \subseteq S$.

Example 9.2. Consider the binary relations R, S on \mathbb{N} defined by

$$\begin{aligned} R &= \{\langle n, m \rangle \mid n < m\} \\ S &= \{\langle n, m \rangle \mid n \leq m\} \end{aligned}$$

Then S extends R .

Example 9.3. Consider the partial order \prec on $\{a, b, c, d, e\}$ which is

$$\{\langle d, b \rangle, \langle d, a \rangle, \langle d, e \rangle, \langle d, c \rangle, \langle a, b \rangle, \langle e, b \rangle, \langle c, b \rangle\}.$$

Then we can extend \prec to the linear order $<$ defined by the transitive closure of

$$d < e < c < a < b.$$

Exercise 9.4. If $\langle A, \prec \rangle$ is a finite partial order, then we can extend \prec to a linear ordering $<$ of A .

Question 9.5. Does the analogous result hold if $\langle A, \prec \rangle$ is an infinite partial order?

Definition 9.6. If A is a set and $n \geq 1$, then

$$A^n = \{\langle a_1, \dots, a_n \rangle \mid a_1, \dots, a_n \in A\}.$$

An n -ary relation on A is a subset $R \subseteq A^n$.

An n -ary operation on A is a function $f: A^n \rightarrow A$.