

Exercise: Use computer to draw a more accurate picture (or use compass and straightedge.)

Lemma:  $\text{Aut}(\mathcal{F}) \cong \text{PGL}(2, \mathbb{Z})$   
 $[GL(2, \mathbb{Z}) / \pm \text{Id}]$

Pf:  $GL \rightarrow \text{Aut}(\mathcal{F})$  }  $A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$   
 $\psi$  }  $f_A(z) = \frac{az+bp}{c_0+dp}$   
 $A \mapsto f_A$

Since  $\det A \begin{pmatrix} r & r \\ p & s \end{pmatrix} = \pm \det \begin{pmatrix} p & r \\ r & s \end{pmatrix}$   
 this is well defined. Notice:  $\pm \text{Id} \in \text{kernel}$ .

Exercise:  $\pm \text{Id} = \text{kernel}$ . Now we must prove surjectivity. [will spend time on this b/c proof is a model for pf of Iverm.]

Step 1: fix  $f \in \text{Aut}(\mathcal{F})$ . Suppose  $f(\infty) = \frac{p}{q}$   
 choose  $r, s$  s.t.  $A = \begin{pmatrix} s & p \\ r & q \end{pmatrix} \in GL(2, \mathbb{Z})$   
 [possible because  $qs - rp = 1$ ]  
 s.  $f_{A^{-1}} \circ f(\infty) = \infty$ . So  $\infty$  now fixed.

Note Step 2:  $f = f_{\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}}$  fixes  $\infty$   
 and  $f = f_{\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}}$  [check]

s.  $\exists h \in \mathbb{Z}$  s.t.  $f_{\begin{pmatrix} 1 & h \\ 0 & 1 \end{pmatrix}} \circ f$  fixes  $\infty$  and 0.

This because  $\{0, \infty\} \in \mathcal{F} \Rightarrow f_{A^{-1}} \circ f$  is adjacent to  $\infty$ , i.e.  $f_{A^{-1}} \circ f(0) \in \mathbb{Z}$ .

Step 3:  $f_{\begin{pmatrix} 1 & h \\ 0 & 1 \end{pmatrix}} \circ f_{A^{-1}} \circ f$  fixes  $\{\pm 1\}$

Let  $g = \text{Id}$  or  $f_{\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}}$  so that

$\phi = g \circ f_{\begin{pmatrix} 1 & h \\ 0 & 1 \end{pmatrix}} \circ f_{A^{-1}} \circ f$  fixes  $\{-1, 0, \infty\}$  pairwise.

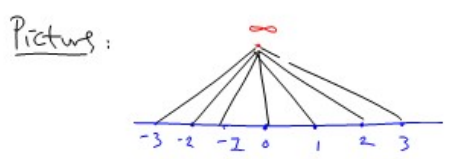
Def:  $X$  a simplicial complex,  $\sigma \in X$  a simplex  
 $lk(\sigma) = \text{link of } \sigma$   
 $= \{ \tau \in X \mid \sigma \cup \tau \in X, \tau \cap \sigma = \emptyset \}$

Picture:  $lk(\infty) = \mathbb{Z}$

Notice:  $\text{Aut}(\mathbb{Z}) \cong \mathbb{Z} \rtimes \mathbb{Z}_2$   
 [infinite dihedral gr.]

WTS  $\phi = \text{Id} \in \text{Aut}(\mathcal{F})$ .

Step 4a:  $\phi|_{\infty \cup lk(\infty)} = \text{Id}$

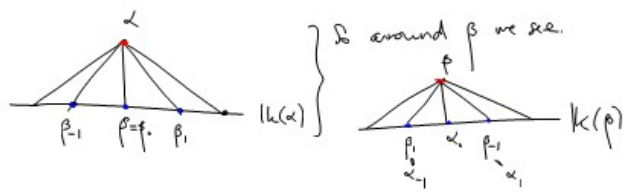


Since  $\phi \in \text{Aut}(\mathcal{F})$ ,  $\phi(\infty) = \infty \Rightarrow \phi|_{lk(\infty)}$   
 is in  $\text{Aut}(\mathbb{Z})$ .  $\therefore \left. \begin{matrix} \phi(0) = 0 \\ \phi(-1) = -1 \end{matrix} \right\} \Rightarrow \phi|_{lk(\infty)} = \text{Id}$ .

**Step 4b** If  $\{\alpha, \beta\} \in \mathcal{F}$  and

$\phi|_{\text{pull}(\alpha)} = \text{Id}$  then  $\phi|_{\text{pull}(\beta)} = \text{Id}$

[Crawling] Pf of this: by picture:



So:  $\phi$  fixes  $\beta$  b/c  $\beta \in \text{lk}(\alpha)$ .  
 $\Rightarrow \phi$  preserves  $\text{lk}(\beta)$ . But  $\phi$  fixes  $\alpha_i, \alpha_{i-1}$   
 $\Rightarrow \phi|_{\text{lk}(\beta)} = \text{Id}$ .

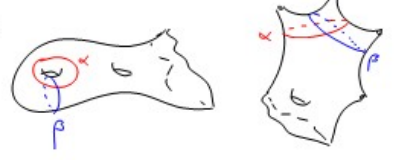
**Step 5** By induction on  $d_{\mathcal{F}}(\infty, p_i)$    
 $\phi(p_i) = p_i$ . Use step 4b. // Lemma

[Here  $d_{\mathcal{F}}(\infty, p_i)$  is distance as measured in  $\mathcal{F}^1$  the one-skeleton.]

**Def:** If  $\alpha, \beta$  nmscp and  $c(\alpha, \beta) = 1$  say  $\alpha, \beta$  dual.

If  $\alpha, \beta$  both pants curves and  $c(\alpha, \beta) = 2$  say  $\alpha, \beta$  dual.

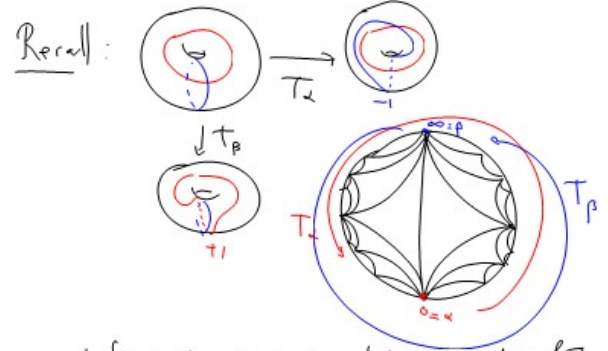
Pictures:



**Def:**  $\alpha \subset S$  is a pants curve if  $\alpha$  separating and  $S_\alpha$  has a  $S_{0,3}$  component.

**Def:** for  $S = S_{1,1}, S_1, S_{0,4}$  define  $\mathcal{C}(S)$  to have vertex set  $= \mathcal{P}(S)$  and simplices given by duality.

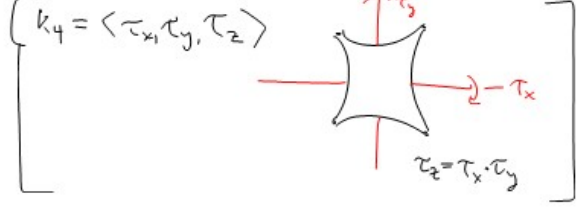
**Ex:**  $\mathcal{C}(S) \cong \mathcal{F}$  in all three cases.



Left twists give counter-clockwise parallelisms of  $\mathcal{F}$

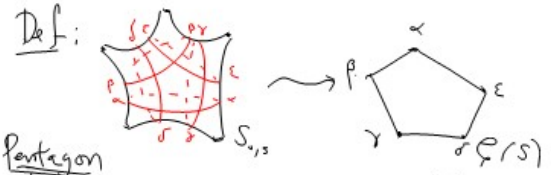


**Lemma:**  $\mathcal{M}(\mathcal{C}(S_{0,4})) \cong K_4 \rtimes \text{PGL}(2, \mathbb{Z})$



**Pf:** Exercise [Hint: Look for short exact sequence that splits.]

**Next time:** Ivanov's Theorem in genus zero.



**Lemma:** The pentagon is unique. Pf Next time